

Discrete Mathematics

Quiz # 4 (April, 2015)

Name: _____ ID: _____

1. (30%) Using strong induction to prove that every positive integer n can be written as a sum of distinct powers of two, that is, as a sum of a subset of the integers $2^0, 2^1, 2^2, \dots, 2^n \dots$ and so on.

Let $P(n)$ be the statement “ n can be written as a sum of distinct powers of two”

Basis step: $P(1)$ is true because $2^0 = 1$

Inductive step:

Assume $P(i)$ is true $\forall i \leq k, i \in \mathbb{N}$.

If $k+1$ is even, $\frac{k+1}{2}$ is an integer smaller than k , and $P\left(\frac{k+1}{2}\right)$ is true. Therefore,

$P(k+1)$ is also true because $k+1$ can be written as $2 * \left(\text{sum of } \frac{k+1}{2}\right)$.

If $k+1$ is odd, it can be written as $2^0 + 2 * \left(\text{sum of } \frac{k}{2}\right)$, which means $P(k+1)$ is true.

This completes the inductive step.

2. (30%) For a full binary tree T , $n(T)$ and $h(T)$ denote the number of vertices and the height of T , respectively. $h(T) = 0$ when T consists of only one vertex (i.e, root). Show that $n(T) = 2^{h(T)+1} - 1$ by using structural induction when the set of full binary trees is defined recursively by the following two steps:

Basis step: A single vertex is a full binary tree.

Recursive step: If T_1 and T_2 are disjoint full binary trees with $h(T_1) = h(T_2)$, $T_1 \cdot T_2$ is also a full binary tree, where $T_1 \cdot T_2$ represents a binary tree consisting of a root r together with the edges connecting r to each of the roots of the left subtree T_1 and the right subtree T_2 .

Basis step:

For the full binary tree consisting of just a root, the result is true because $n(T)=1$ and $h(T)=0$, and $1 = 2^{0+1} - 1$.

Inductive step:

Assume that $n(T_1) = 2^{h(T_1)+1} - 1$, $n(T_2) = 2^{h(T_2)+1} - 1$ where $h(T_1) = h(T_2)$.

**Then $n(T) = 1 + n(T_1) + n(T_2) = 1 + 2^{h(T_1)+1} - 1 + 2^{h(T_2)+1} - 1$
 $= 2^{h(T_1)+1+1} - 1 = 2^{h(T)+1} - 1$**

3. (40%) Use a recursive algorithm to represent each of the following functions.

(a) $f(n, m) = n! \bmod m$, where $n, m \in \mathbb{N}$.

(b) $g(a, n) = a^{2^n}$, where $a \in \mathbb{R}$ and $n \in \mathbb{N}$.

(a) Procedure modfactorial(n, m : positive integers)

If $n=1$ then modfactorial(n, m):=1

Else modfactorial(n, m):= (n *modfactorial($n-1, m$)) mod m

(b) Procedure a(n : positive integer)

If $n=1$ then $a(n) := a^2$

Else $a(n) := (a(n-1))^2$